

A note on the characterization of optimal allocations in OLG economies with multiple goods

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- 1 Introduction
- 2 The model
- 3 Characterization of Pareto optimal allocations
- 4 Sketch of the proof

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Introduction

A very standard question in optimization: first-order characterization of Pareto optimal allocation;

A non-standard framework: countably many objective functions;

A non-standard space: $\prod_{t=1}^{\infty} \mathbb{R}^{L_t}$ and the product topology.

A direct proof based on Balasko-Shell

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A standard OLG exchange economy \mathcal{E}

Infinitely many dates ($t = 1, 2, \dots$),

at date t ,

- a finite set \mathcal{L}_t of available commodities;
- a finite and non-empty set of consumers \mathcal{I}_t living two periods, consumption set $X^i = \mathbb{R}_+^{L_t} \times \mathbb{R}_+^{L_{t+1}}$, endowments $e^i \in \mathbb{R}_{++}^{L_t} \times \mathbb{R}_{++}^{L_{t+1}}$;
- for each consumer i , a strict preference relation P_i from X^i to X^i .

Notations: $\bar{e}_t = \sum_{i \in \mathcal{I}_{t-1} \cup \mathcal{I}_t} e_t^i$

Assumption A.

- a) For all i in \mathcal{I} , $P^i(x^i)$ open in X , convex,
 $x^i \in \bar{P}^i(x^i)$ and $x^i \notin P^i(x^i)$,
 $P^i(x^i) + (\mathbb{R}_+^{L_t} \times \mathbb{R}_+^{L_{t+1}}) \subset P^i(x^i)$.
- b) For all $i \in \mathcal{I}$, for all x_i in the interior of X_i ,
 $-N_{\bar{P}^i(x^i)}(x^i)$ is a half line $\{t\gamma^i(x^i) \mid t \geq 0\}$,
 $\gamma^i(x^i) \in \mathbb{R}_{++}^{L_t} \times \mathbb{R}_{++}^{L_{t+1}}$ continuous mapping on the
interior of X^i , $\|\gamma^i(x^i)\| = 1$.

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Pareto optimal allocations

Definition

x feasible is Pareto optimal (PO) (resp. weakly Pareto optimal (WPO)) if there is no (y^i) in $\prod_{i \in \mathcal{I}} X^i$ such that:

$$\sum_{i \in \mathcal{I}_{t-1} \cup \mathcal{I}_t} y_t^i = \bar{e}_t, \text{ for } t \geq 1$$

and for all $i \in \mathcal{I}$, $y^i \in \bar{P}^i(x^i)$, with $y^i \in P^i(x^i)$ for at least one individual i (resp. $\exists \bar{t}$ such that $\forall t \geq \bar{t}, \forall i \in \mathcal{I}_t y^i = x^i$).

Characterization of WEAK Pareto optimal allocations

Definition

A feasible allocation $x = (x^i)$ is supported by the price p if for each $t \in \mathbb{N}$, for all $i \in \mathcal{I}_t$ and for all $\xi^i \in P^i(x^i)$, $\Pi_t \cdot \xi^i > \Pi_t \cdot x^i$, where $\Pi_t = (p_t, p_{t+1})$.

Lemma

The feasible interior allocation $x = (x^i)$ is WPO if and only if there exists a price sequence p which supports $x = (x^i)$. Then Π_t is collinear to $\gamma^i(x^i)$ for all t and $i \in \mathcal{I}_t$.

Additional assumptions

$x = (x^i) \in \prod_{i \in \mathcal{I}} X^i$ a feasible allocation: **Assumption C:**

$\exists \underline{r} > 0, \forall i \in \mathcal{I}, \bar{B}_{t,t+1}(x^i + \underline{r}\gamma^i(x^i), \underline{r}) \setminus \{x^i\} \subset P^i(x^i);$

Assumption C': $\exists \bar{r} > 0, \text{ for all } i \in \mathcal{I}, \forall \xi^i \leq (\bar{e}_t, \bar{e}_{t+1}) \text{ if } \xi^i \in P^i(x^i), \text{ then } \xi^i \in \bar{B}(x^i + \bar{r}\gamma^i(x^i), \bar{r});$

Assumption G: $\exists \bar{\eta} > \underline{\eta} > 0, \forall i \in \mathcal{I}_{-0}, \underline{\eta} \leq \frac{\|\gamma_t^i(x^i)\|}{\|\gamma_{t+1}^i(x^i)\|} \leq \bar{\eta};$

Assumption B: $\exists \bar{e} \in \mathbb{R}_{++}^L, \exists \underline{\varepsilon} > 0, \forall t \in \mathbb{N}^*, \bar{e}_t \leq \bar{e} \text{ and } \forall i \in \mathcal{I}_{-0}, (\underline{\varepsilon}\mathbf{1}_t, \underline{\varepsilon}\mathbf{1}_{t+1}) \leq x^i.$

Proposition

$x = (x^i) \in \prod_{i \in \mathcal{I}} X^i$ an allocation supported by
 $p = (p_1, p_2, \dots, p_t, \dots)$ satisfying Assumptions A, C, C', G and B.
Then, x is Pareto optimal if and only if:

$$\sum_{t \in \mathbb{N}^*} \frac{1}{\|p_t\|} = +\infty.$$

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Aggregation: one consumer at each generation

$$\bar{P}_t((x^i)) := \sum_{i \in \mathcal{I}_t} \bar{P}^i(x^i)$$

$$\bar{x}^t := \sum_{i \in \mathcal{I}_t} x^i$$

$$N_{\bar{P}_t((x^i))}(\bar{x}^t) = \bigcap_{i \in \mathcal{I}_t} N_{\bar{P}^i(x^i)}(x^i)$$

x Pareto optimal if and only if \bar{x} Pareto optimal for the preference relation \bar{P}_t . x supported by a price, then \bar{x} supported by the same price.

A geometric approach

Feasible aggregate transfer: $\bar{h} = (\bar{h}_t^t, \bar{h}_{t+1}^t)$ with $\bar{h}_{t+1}^t = -\bar{h}_{t+1}^{t+1}$.

Net present value: $\eta^t = \Pi_t \cdot \bar{h}^t$

$$\alpha^t = \frac{\|\bar{h}^t\|^2 \|\Pi_t\|}{\eta^t}$$

Lemma

Under Assumption C', \bar{h} Pareto improving implies α is bounded above.

Lemma

\bar{h} feasible Pareto improving aggregate transfer, then:

$$\begin{aligned} \alpha^t &\geq \frac{\|\Pi_t\|}{\eta^t} \|\bar{h}^t\|^2 \\ &= \frac{\|\Pi_t\|}{\eta^t} \left[\frac{1}{\|p_t\|^2} (\eta^0 + \dots + \eta^{t-1})^2 + \frac{1}{\|p_{t+1}\|^2} (\eta^0 + \dots + \eta^t)^2 \right] \end{aligned}$$

Lemma

η a positive sequence in \mathbb{R} .

$$\bar{h}_{t+1}^t = \left(\eta^0 + \eta^1 + \dots + \eta^t \right) \frac{p_{t+1}}{\|p_{t+1}\|^2} \text{ and } \bar{h}_t^t = -\bar{h}_t^{t-1}$$

$$\begin{aligned} \alpha^t &= \frac{\|\Pi_t\|}{\eta^t} \|\bar{h}^t\|^2 \\ &= \frac{\|\Pi_t\|}{\eta^t} \left[\frac{1}{\|p_t\|^2} (\eta^0 + \dots + \eta^{t-1})^2 + \frac{1}{\|p_{t+1}\|^2} (\eta^0 + \dots + \eta^t)^2 \right] \end{aligned}$$

Under Assumptions B and C, if α bounded then $\exists \mu > 0$ such $\mu \bar{h}$ is a feasible Pareto improving aggregate transfer.